

Roll Your Own Test Bed for Embedded Real-Time Protocols: A Haskell Experience

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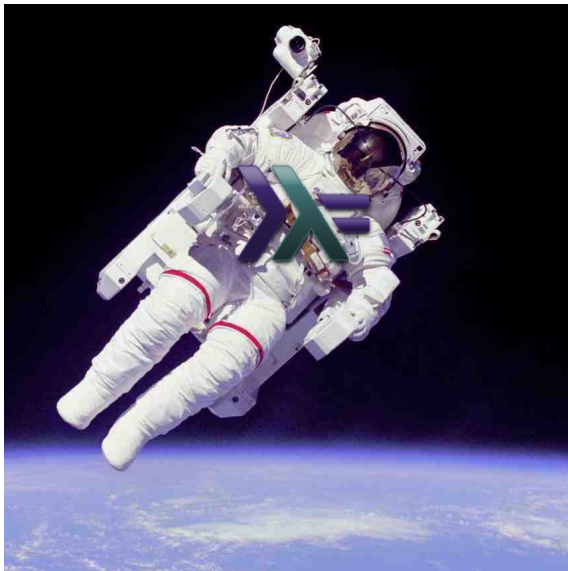
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This is a story about using Haskell. . .



...Haskell in Space!



Ok, just kidding.

Goals:

1. Teach you enough about **physical-layer protocols** to make you dangerous.
2. Tell you how I easily modeled **real-time distributed systems** in a lazy, pure language.
3. Tell you how I used QuickCheck as a “**probability calculator**”.

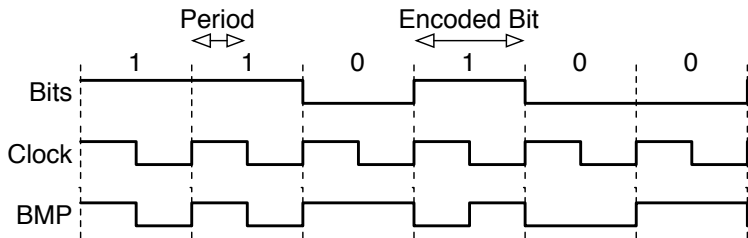
(1) Physical-Layer Protocols

Do you wonder how

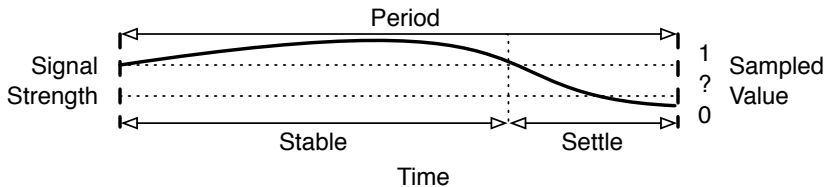
- ▶ The credit-card slider reads your magnetic stripe?
- ▶ The CD player reads your Rolling Stones CD?
- ▶ The internet gets transmitted to your computer?

All are achieved using physical-layer protocols: a transmitter t_x sends a receiver r_x a bit-stream in real-time.

Biphase Mark Protocol



Signal Strength Over Time



Parameters

```
-- | Realtime input parameters.
data Params = Params
  { tPeriod    :: Time -- ^ Tx's clock period.
  , tSettle    :: Time -- ^ Nominal signal settling time.
  , rScanMin   :: Time -- ^ Rx's min scan duration.
  , rScanMax   :: Time -- ^ Rx's max scan duration.
  , rSampMin   :: Time -- ^ Rx's min sampling duration.
  , rSampMax   :: Time -- ^ Rx's max sampling duration.
  } deriving (Show, Eq)
```

Some Constraints

```
paramsConst :: Params -> Bool
paramsConst p =
    0 < tPeriod p -- tPeriod
  && 0 <= tSettle p -- tSettle
  && tSettle p < tPeriod p -- tSettle
  && 0 < rScanMin p -- rScanMin
  && rScanMin p <= rScanMax p -- rScanMax
  && rScanMax p < tStable -- rScanMax
  && tPeriod p + tSettle p < rSampMin p -- rSampMin
  && rSampMin p <= rSampMax p -- rSampMax
  && rSampMax p < tPeriod p + tStable - rScanMax p
  where tStable = tPeriod p - tSettle p
```

(2) Modeling Real-Time in Haskell

Modeling Real-Time

In a Lazy Functional Language

- ▶ **Question:** How do I model distributed real-time behavior in a functional language?

More precisely, we want to model a **partially-synchronous** real-time system with possibly non-deterministic bounds on asynchrony.

- ▶ **Answer:** The discrete-event simulation folks figured this out a few decades ago. We'll just borrow their ideas.

Claim: many practical real-time systems fit this model.

Let's Take Turns!

or Discrete-Event Simulation

- ▶ Suppose you have participants $\rho_0, \rho_1, \dots, \rho_n$.

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- ▶ Suppose you have participants p_0, p_1, \dots, p_n .
- ▶ Suppose the present time is time t .
- ▶ Each p_i is scheduled to take some action in the future, when its **timeout** is reached.
 - ▶ At time $t + n$ (**synchronous**).
 - ▶ Within $(t + n, t + n + \delta)$ (**partially-synchronous**).

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- ▶ p_i takes its action, and updates its timeout.

Example

$$t = 0$$

ρ_0	ρ_1
0	0 from $[0, 1]$

Example

$$t = 0$$

p_0	p_1
0	0 from $[0, 1]$ \Leftarrow

Example

$$t = 0$$

p_0	p_1
0	0 from $[0, 1]$ \Leftarrow
	7

Example

$$t = 0$$

	p_0	p_1
\Rightarrow	0	0 from $[0, 1]$
		7

Example

$$t = 0$$

	p_0	p_1
\Rightarrow	0	0 from $[0, 1]$
		7
	9.89 from $(8.5, 10)$	

Example

$$t = 7$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	

Example

$$t = 7$$

p_0	p_1
0	0 from $[0, 1]$
	7 \leftarrow
9.89 from $(8.5, 10)$	

Example

$$t = 7$$

p_0	p_1
0	0 from $[0, 1]$
	7 \leftarrow
9.89 from $(8.5, 10)$	
	9

Example

$$t = 9$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9

Example

$$t = 9$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9 \leftarrow

Example

$$t = 9$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9 \leftarrow
	10

Example

$$t = 9.89$$

ρ_0	ρ_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9
	10

Example

$$t = 9.89$$

ρ_0	ρ_1
0	0 from $[0, 1]$
	7
\Rightarrow 9.89 from $(8.5, 10)$	
	9
	10

Example

$$t = 9.89$$

ρ_0	ρ_1
0	0 from $[0, 1]$
	7
\Rightarrow 9.89 from $(8.5, 10)$	
	9
	10
$(10, 12)$	

Example

$$t = 10$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9
	10
$(10, 12)$	

Example

$$t = 10$$

p_0	p_1
0	0 from $[0, 1]$
	7
9.89 from $(8.5, 10)$	
	9
	10 \leftarrow
$(10, 12)$	

(3) QuickCheck

A Tale of Two QuickCheck Uses

In One Slide

- ▶ **# 1 Testing:** we'll feed the model QC-generated real-time parameters satisfying the constraints.
About 100,000 tests-runs per minute on a MacBook.
- ▶ **# 2 Probability Calculating:** QC for use in stochastic testing.
- ▶ For both, we use monadic QuickCheck, since the model itself is within the `IO` monad. (A [small patch](#) is needed to QC.)
- ▶ And we use the super-fast `System.Random.Mersenne` for generating timeouts.
- ▶ But with no optimizations, testing is surprisingly fast!

Conclusions

- ▶ Emulating real-time is real easy in a pure, lazy language.
 - ▶ Generating real-time parameters is quick with QuickCheck.
 - ▶ And QuickCheck can be used for probabilistic reliability analysis.
 - ▶ **Google**: `biphase quickcheck` to get the code & QuickCheck patch.
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Shameless plug: I'm looking for a **summer student** (undergrad or Ph.D.) in 2010 and/or 2011 who'd like to do some hacking & research on a NASA-sponsored project. . .

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